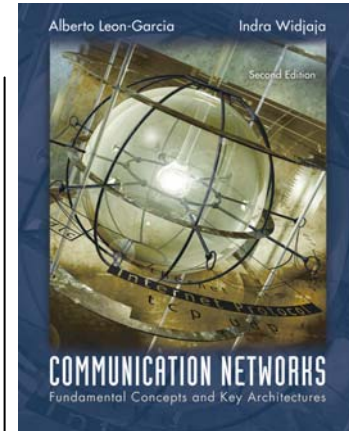


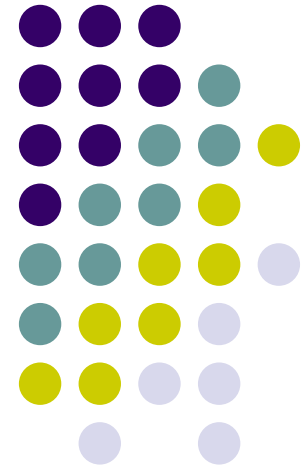
# Chapter 12

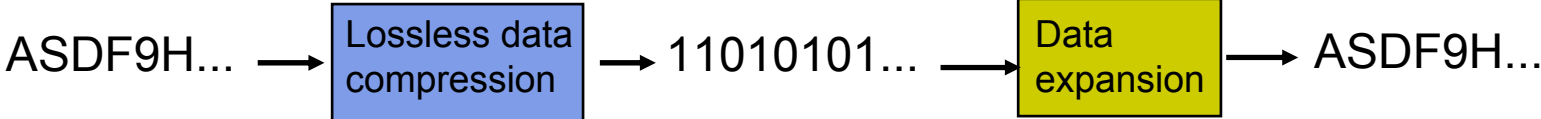
# Multimedia

# Information



## Chapter Figures



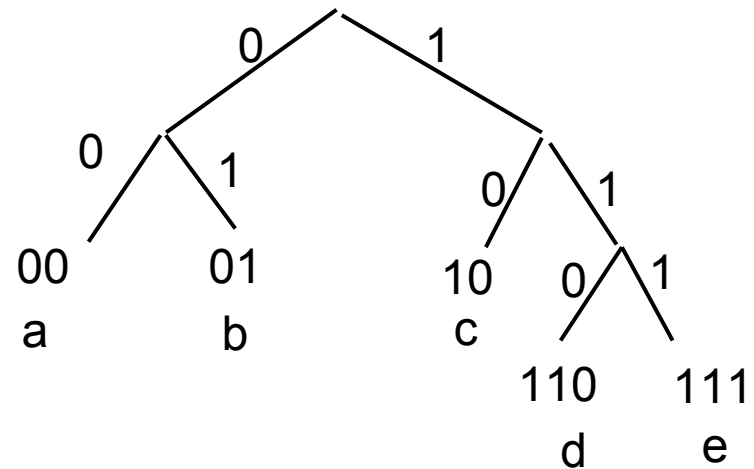


## Assume

- 5 symbol information source: {a,b,c,d,e}
- symbol probabilities: {1/4, 1/4, 1/4, 1/8, 1/8}

Symbol Codeword

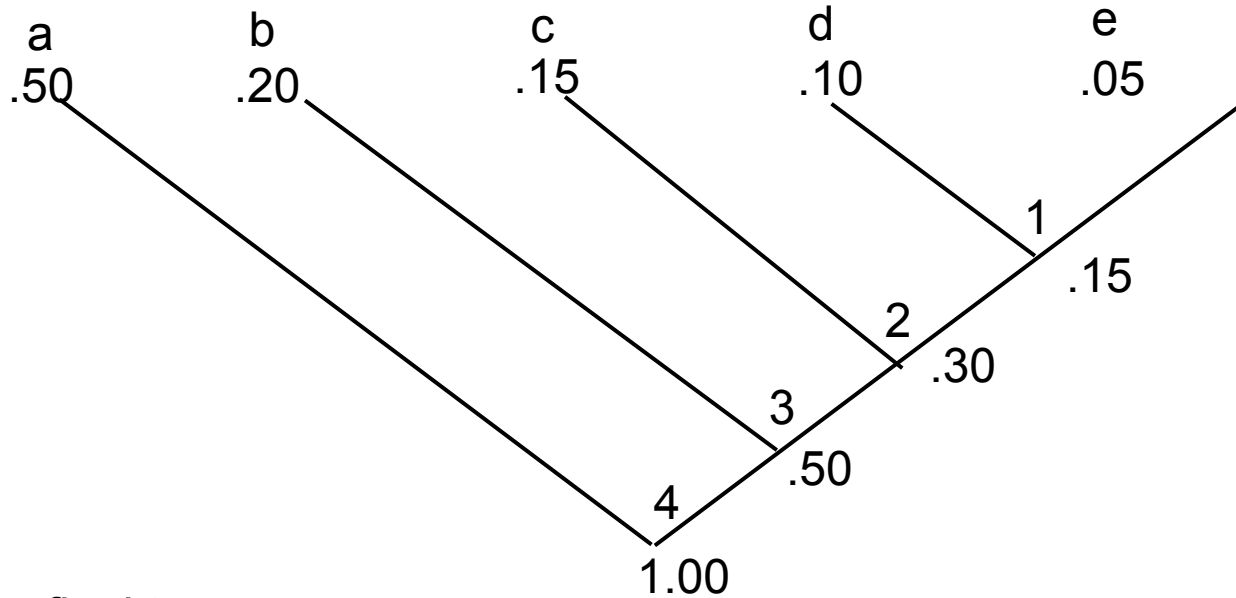
a	00
b	01
c	10
d	110
e	111



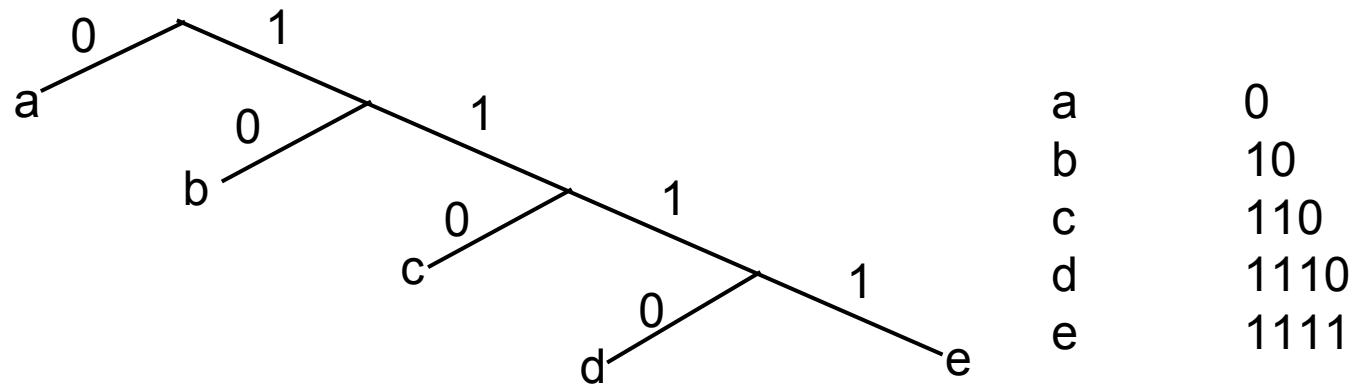
aedbbad.... mapped into 00 111 110 01 01 00 110 ... 17 bits

Note: decoding done without commas or spaces

(a) Building the tree code by Huffman algorithm



(b) The final tree code



- “Blank” in strings of alphanumeric information

-----\$5----3-----\$2-----\$3-----

○ ○ ○ ○ ○ ● ○ ○ ○ ○ ○ ○ ○ ○ ○ ○ ● ○

- “0” (white) and “1” (black) in fax documents

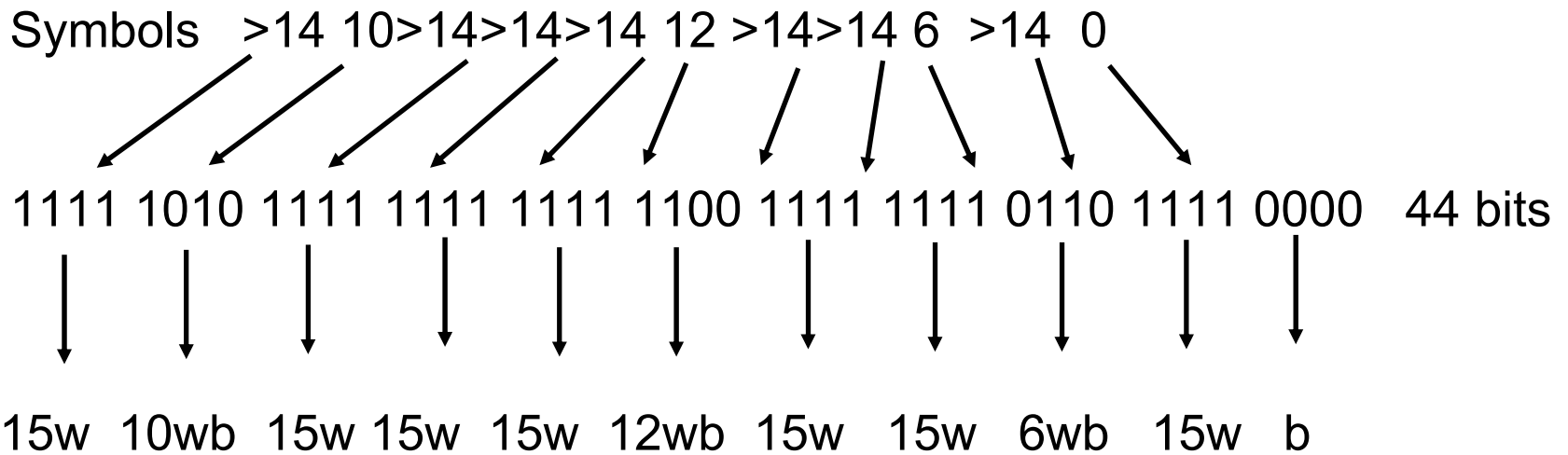
Inputs: ■ □ □ □ □ □ □ □ □ □ □ □ ■ □ □ □ □ □ □ □ □ ■ □ □ □ ■ □ □ □ ■ □ □ □ □ □ □

Run	Length	Codeword	Codeword ( $m = 4$ )
1	0	00..00	0000
01	1	00..01	0001
001	2	00..10	0010
0001	3	00..11	0011
00001	4	.	.
000001	5	.	.
0000001	6	.	.
.	.	.	.
.	.	.	.
000...01	$2^m - 2$	11..10	1110
000...00	run $> 2^m - 2$	11..11	1111

←  $m$  →

Example: Code 1,  $m = 4$

Runs 000...001000...01000...0100...001... 137 bits  
           25          57          36          15



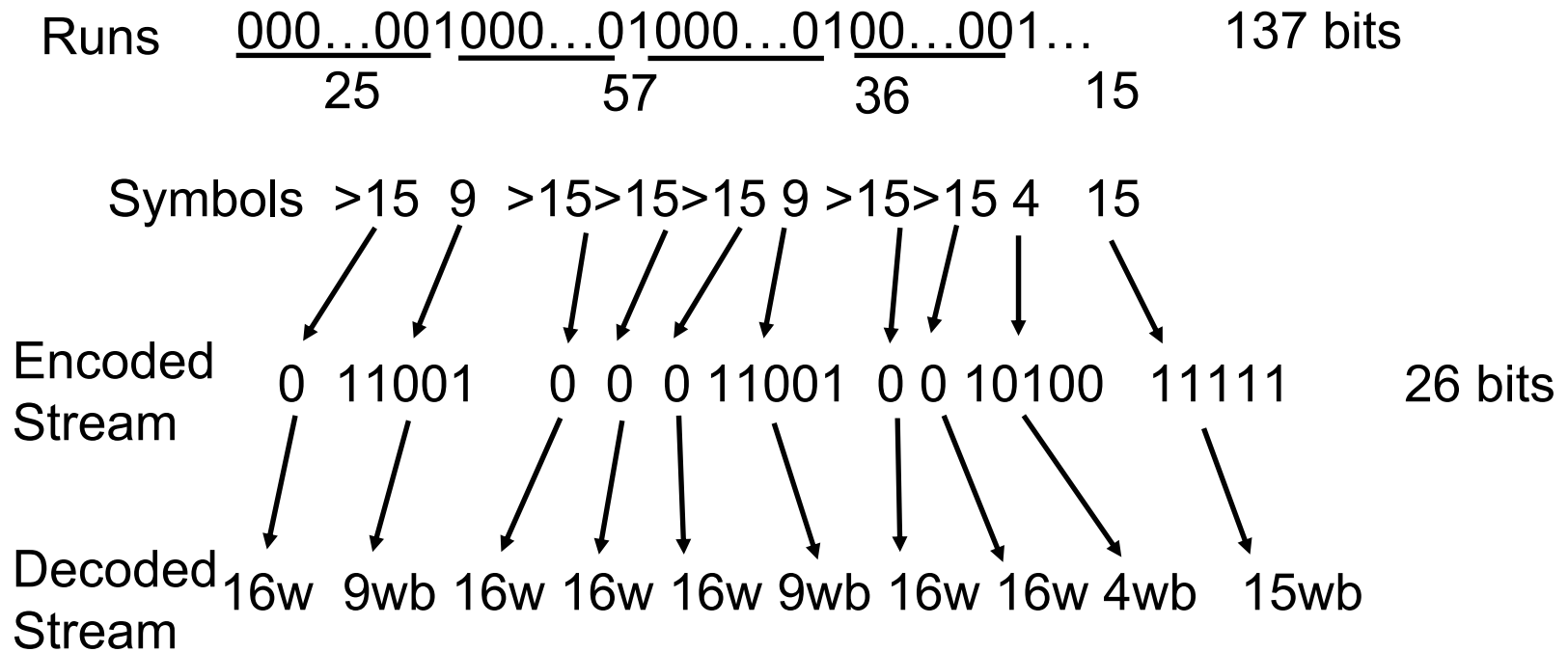
Inputs: ■ □ □ □ □ □ □ □ □ □ □ □ ■ □ □ □ □ □ □ □ □ ■ □ □ □ ■ □ □ □ ■ □ □ □ □ □ □

Run	Length	Codeword	Codeword ( $m = 4$ )
1	0	10..00	10000
01	1	10..01	10001
001	2	10..10	10010
0001	3	10..11	10011
00001	4	.	.
000001	5	.	.
0000001	6	.	.
.	.	.	.
.	.	.	.
000... 01	$2^m - 1$	11..11	11111
000... 00	run $> 2^m - 1$	0	0

←  $m + 1$  →



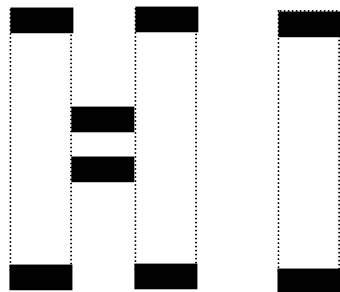
Example: Code 2,  $m = 4$



(a) Huffman code applied to white runs and black runs



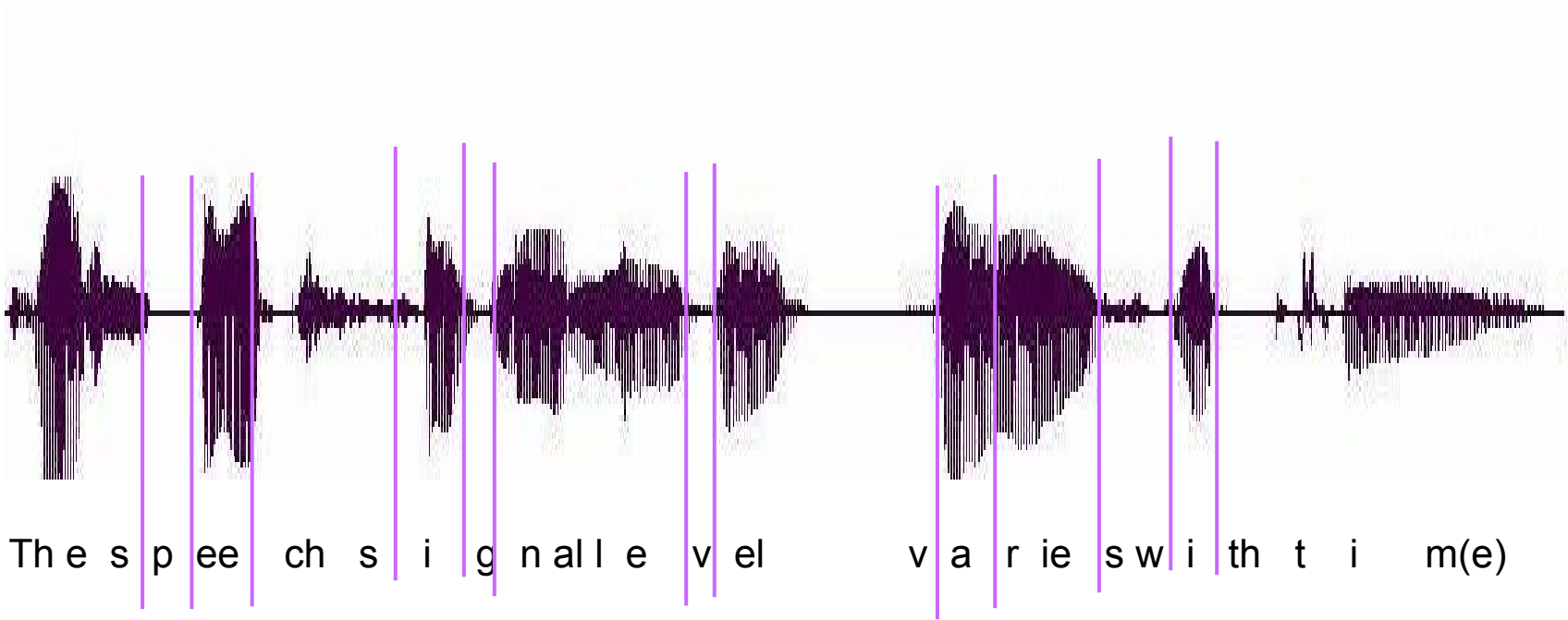
(b) Encode differences between consecutive lines

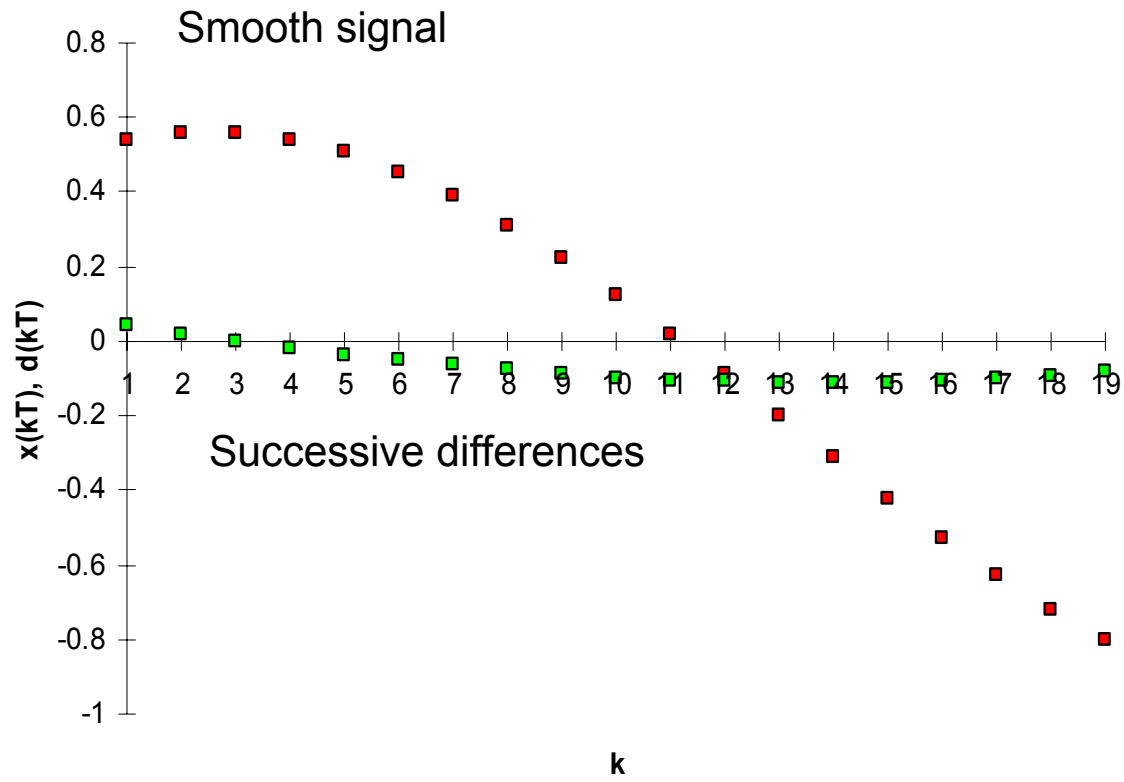


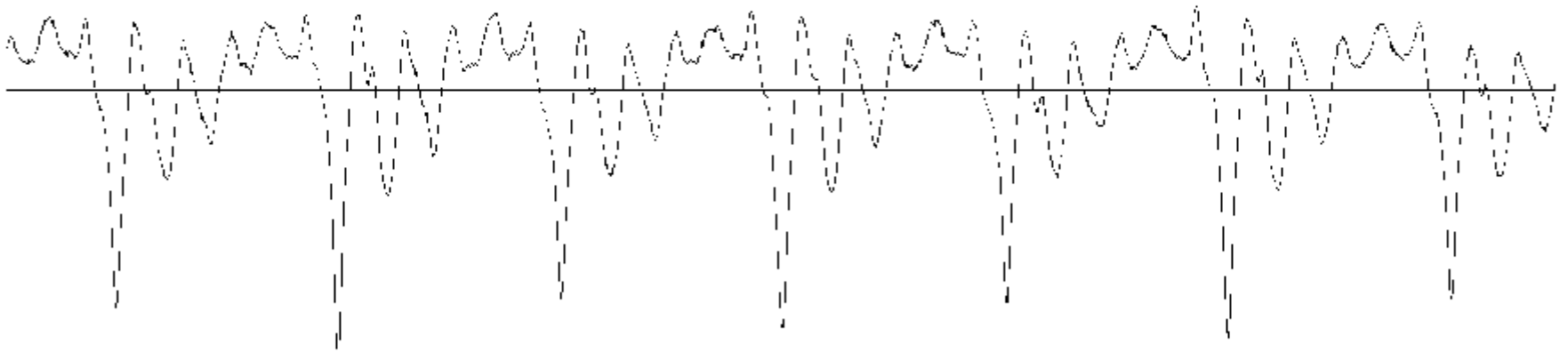
“All tall We all are tall. All small We all are small.”

Can be mapped into:

“All\_ta[2,3]We\_[6,4]are[4,5].\_[1,4]sm[6,15][31,5].”

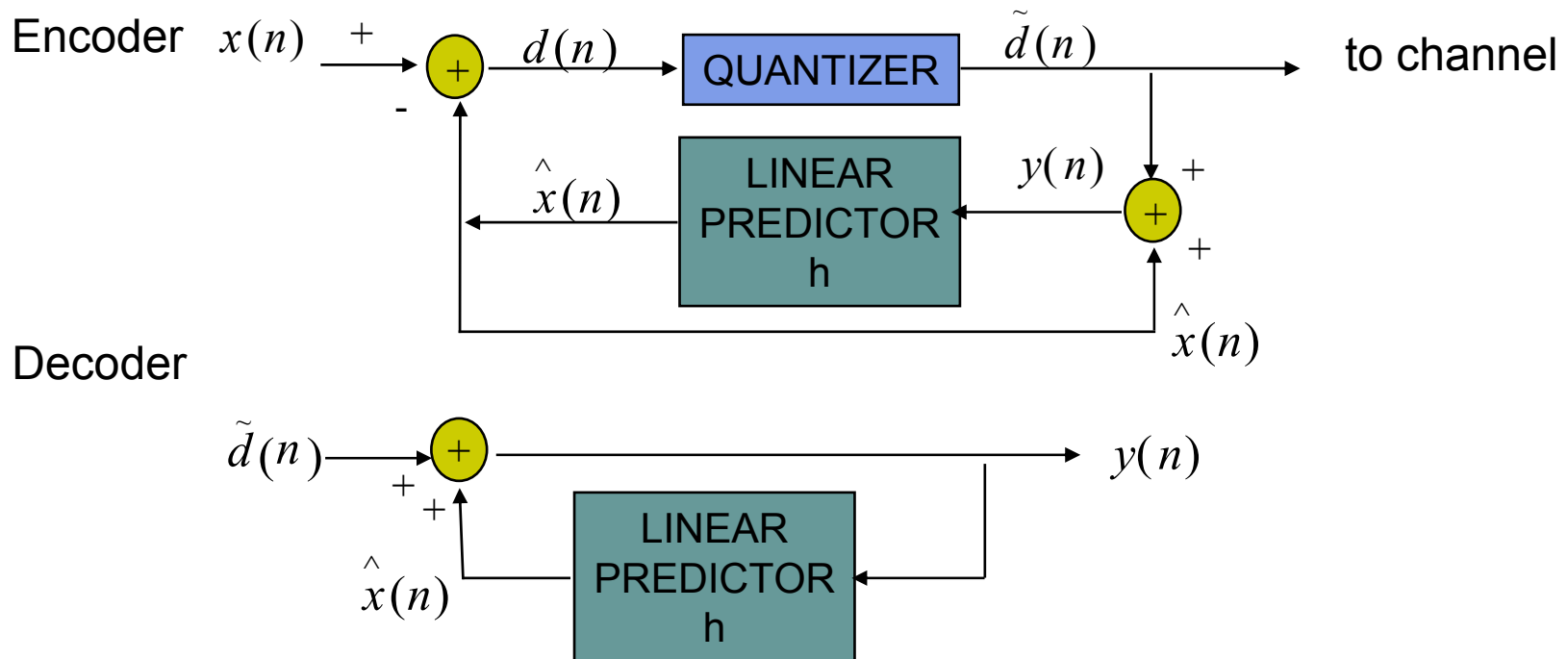








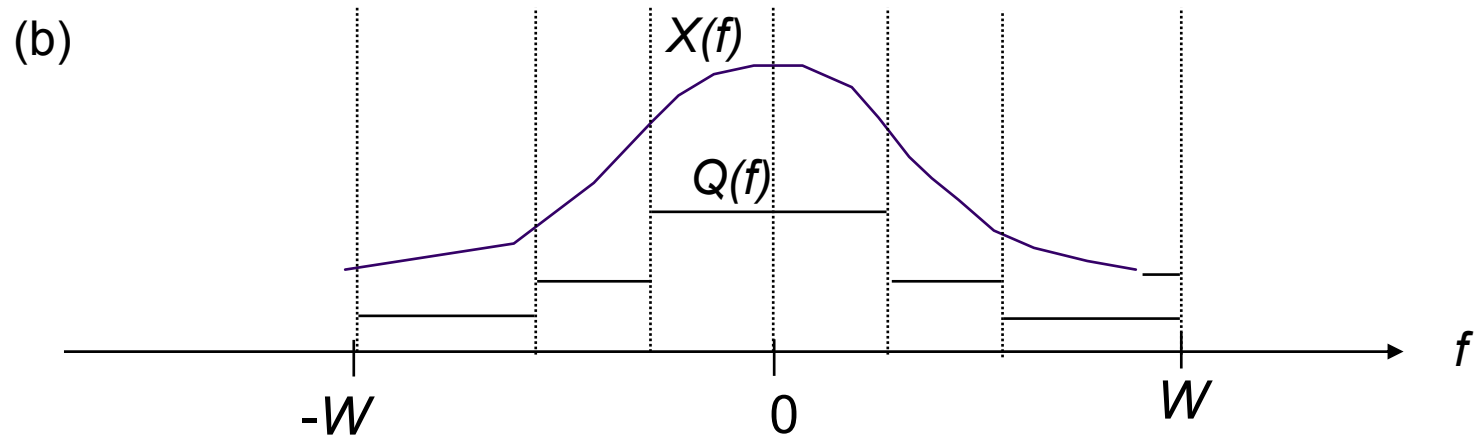
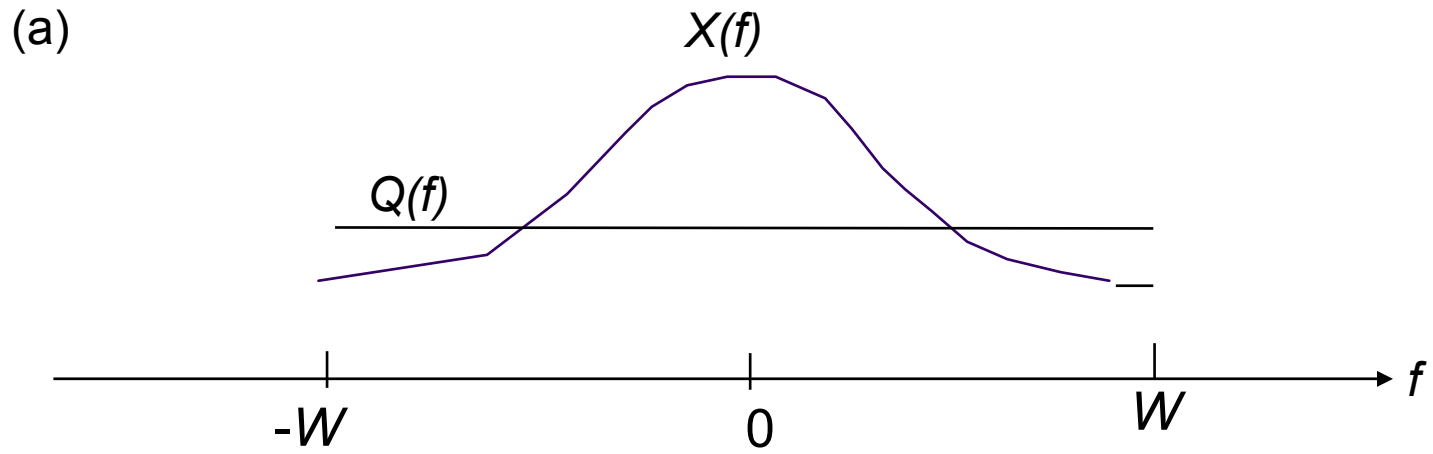
Quantize the difference between prediction and actual signal:



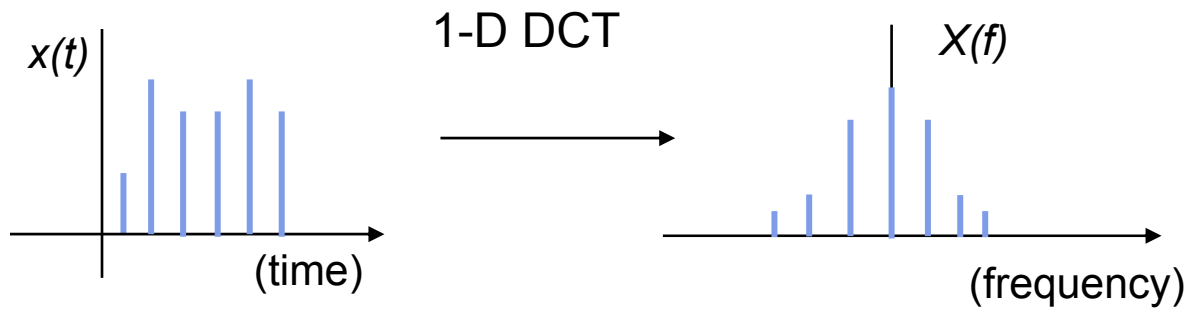
$$y(\tilde{n}) - x(n) = \hat{x}(\tilde{n}) + \tilde{d}(\tilde{n}) - x(n) = \tilde{d}(\tilde{n}) - d(n) = e(n)$$

The end-to-end error is only the error introduced by the quantizer!

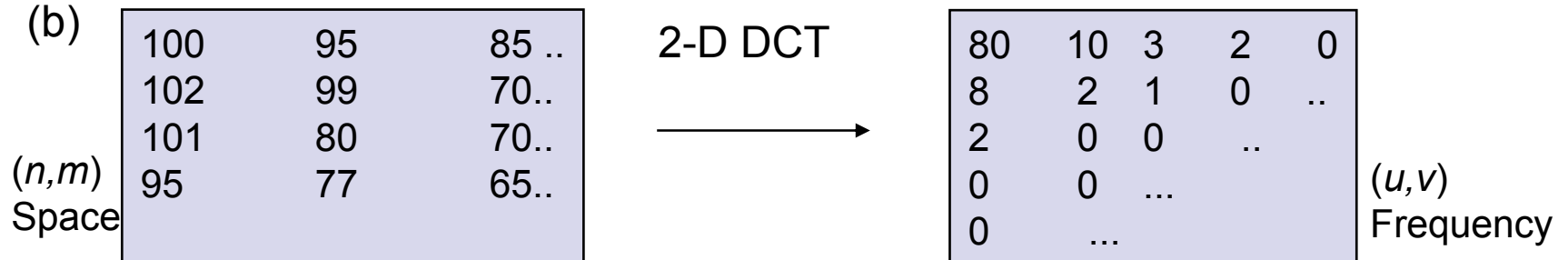


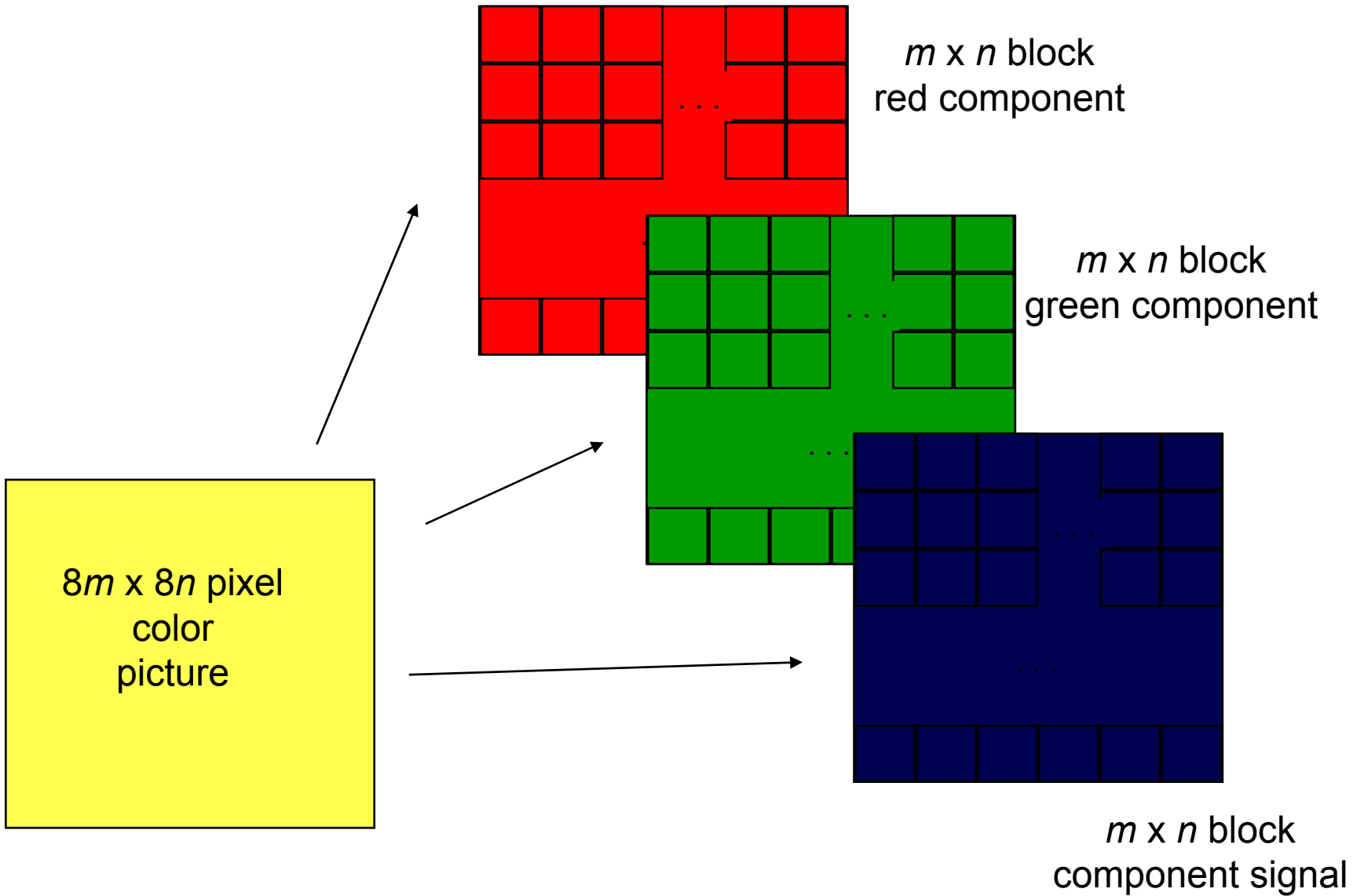


(a)



(b)





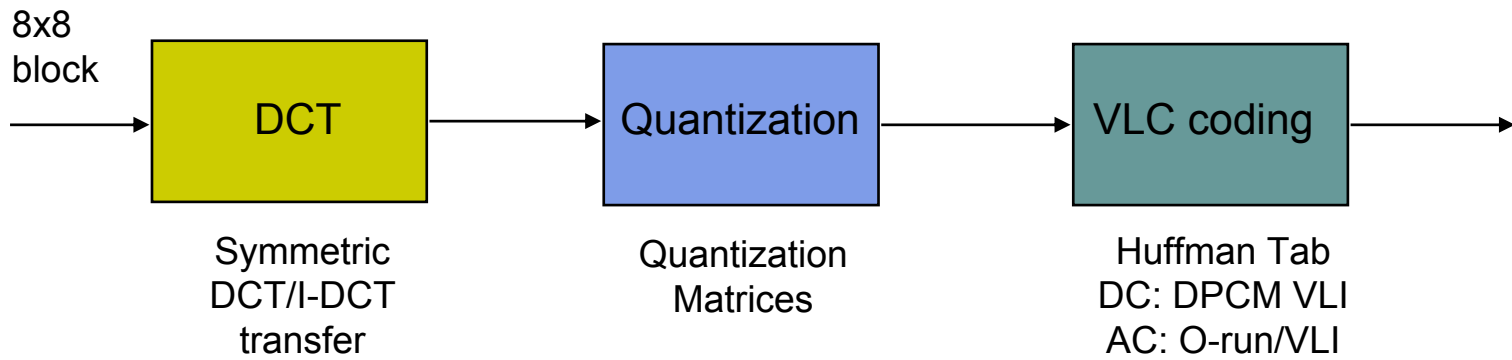
180	150	115	100	100	100	100	100
250	180	128	100	100	100	100	100
190	170	120	100	100	100	100	100
160	130	110	100	100	100	100	100
110	100	100	100	100	100	100	100
100	100	100	100	100	100	100	100
100	100	100	100	100	100	100	100
100	100	100	100	100	100	100	100

8x8 block of 8-bit pixel values

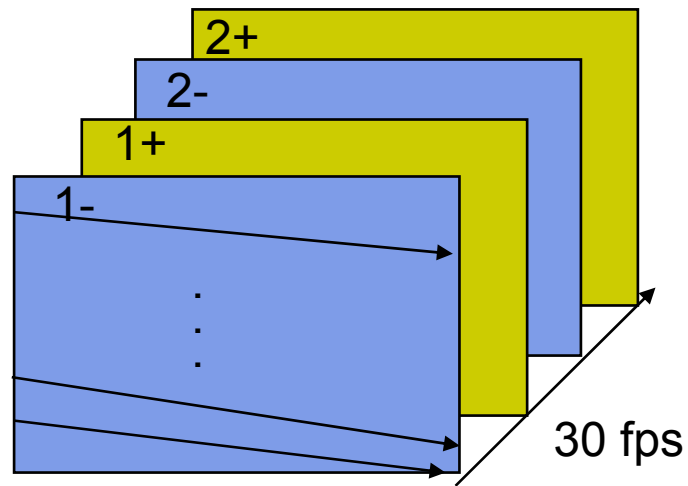
DCT  
→

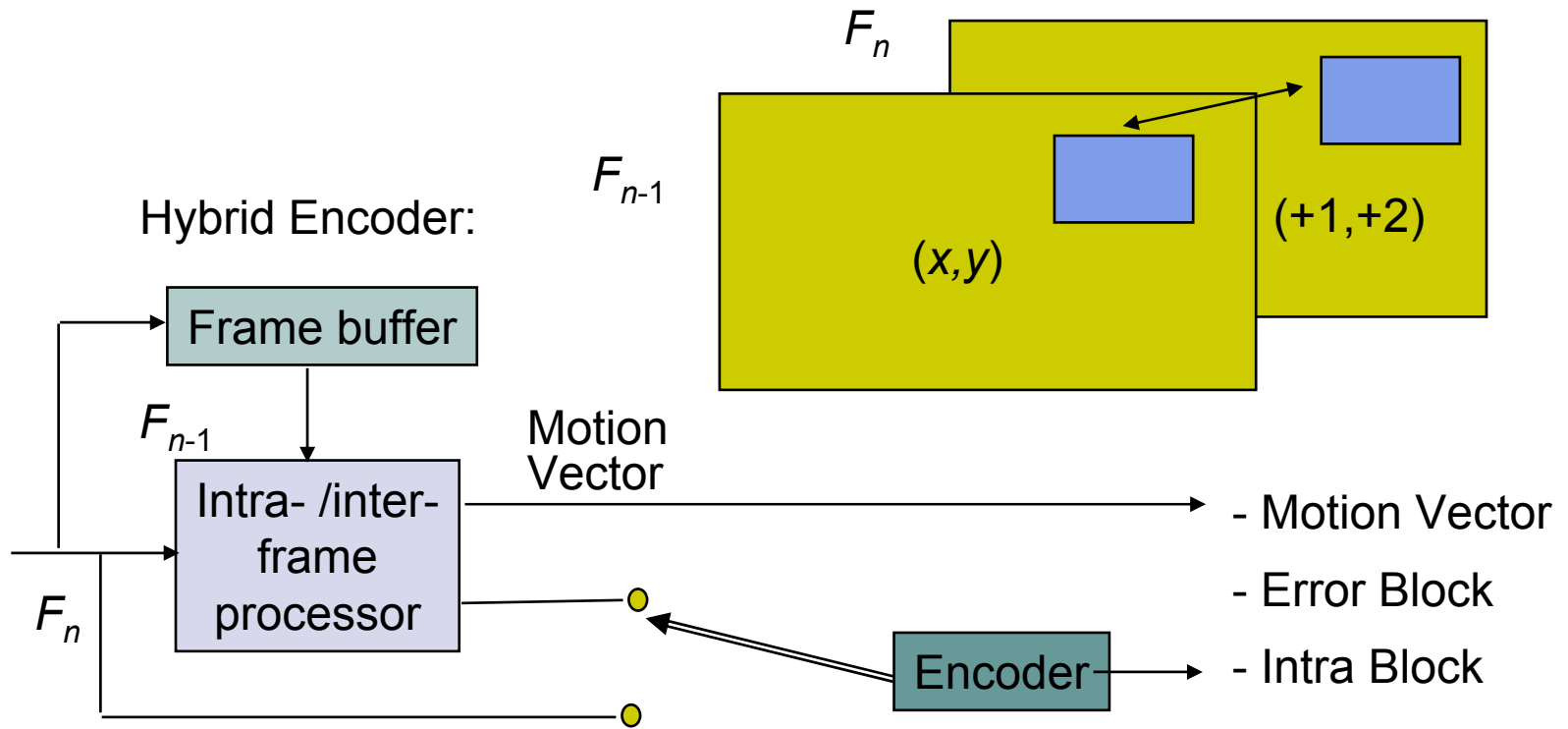
111	22	15	5	1	0	0	0
14	17	10	4	1	0	0	0
2	2	1	0	0	0	0	0
-4	-4	-2	-1	0	0	0	0
-3	-3	-1	0	0	0	0	0
-1	-1	0	0	0	0	0	0
0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0

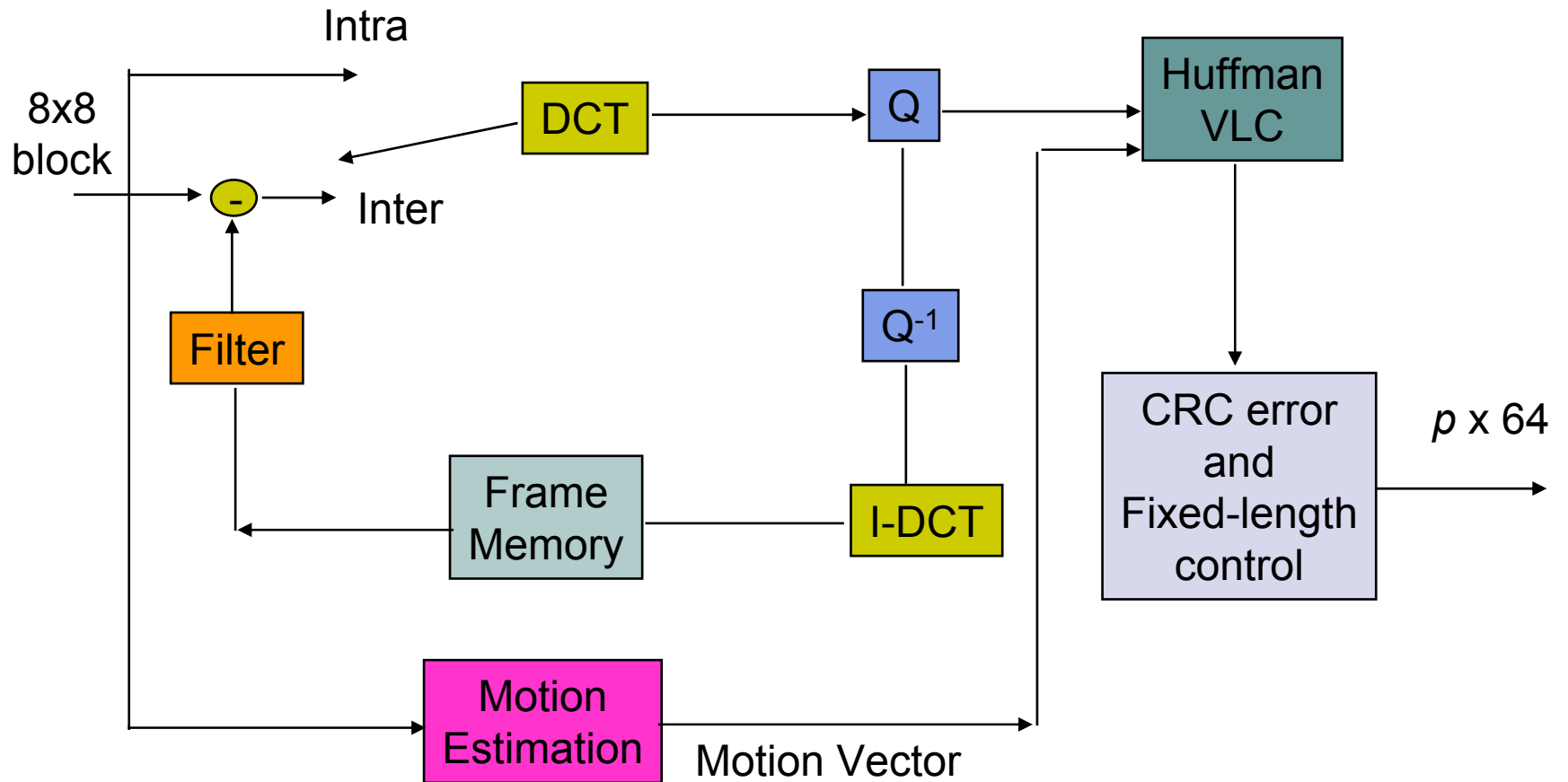
Quantized DCT Coefficients



Info =  $M$  bits/pixel x  $(W \times H)$  pixels/frame x  $F$  frames/second

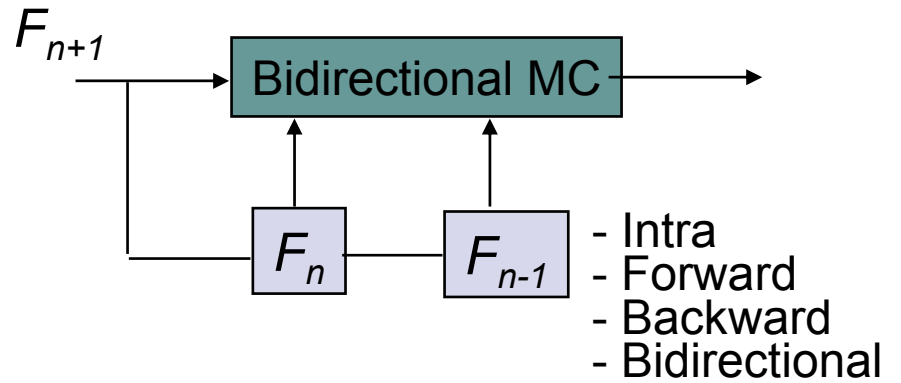
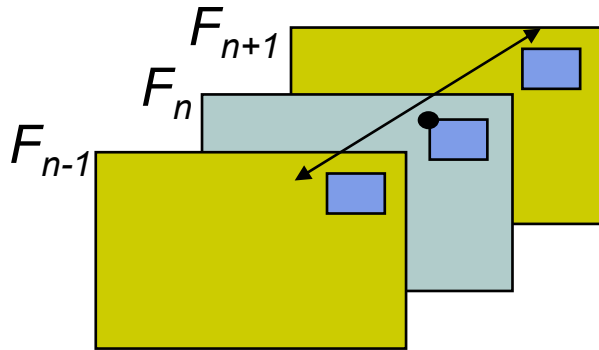
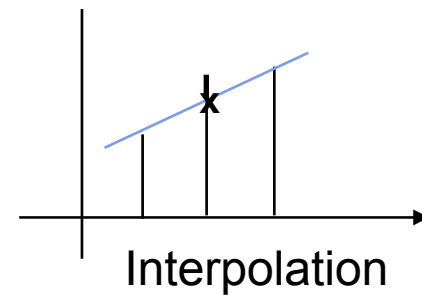
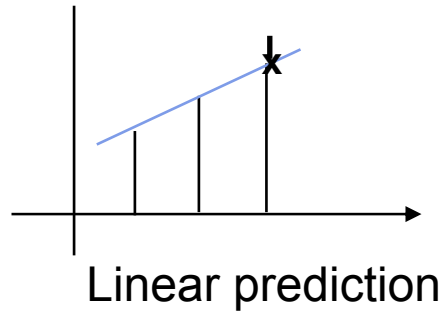
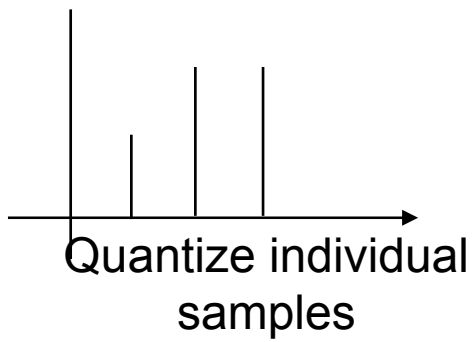




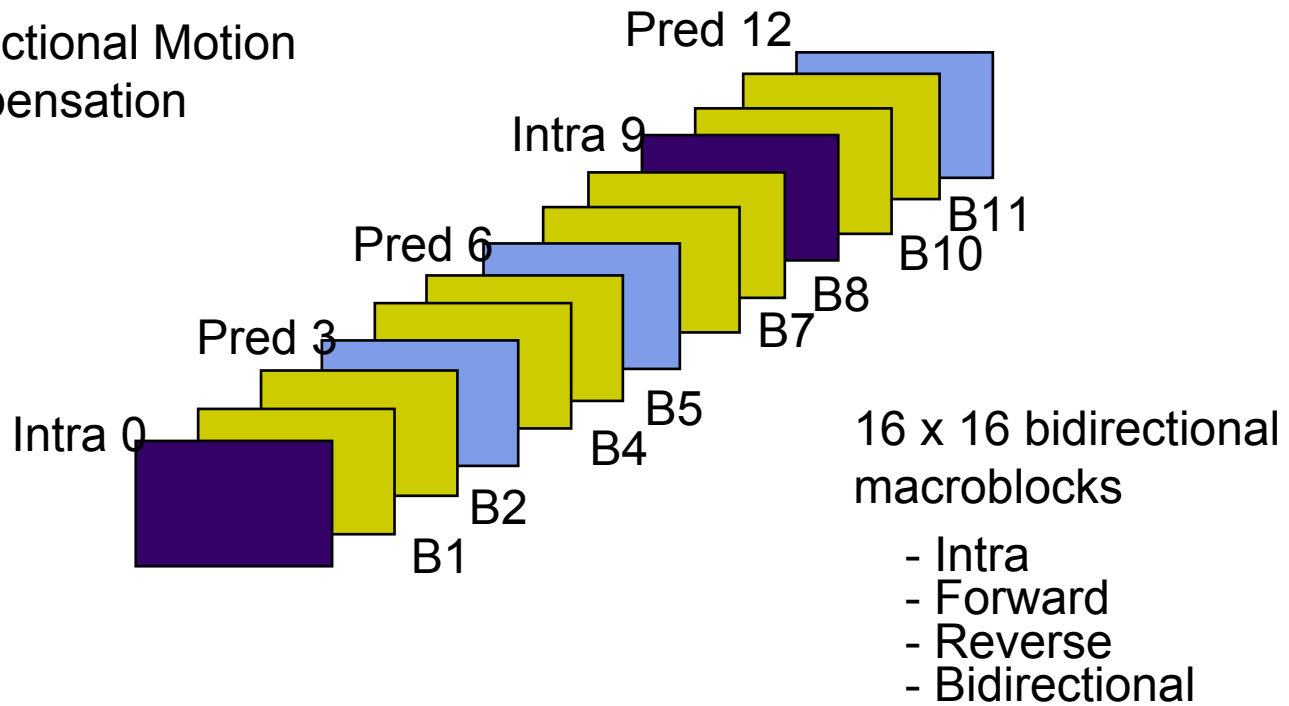




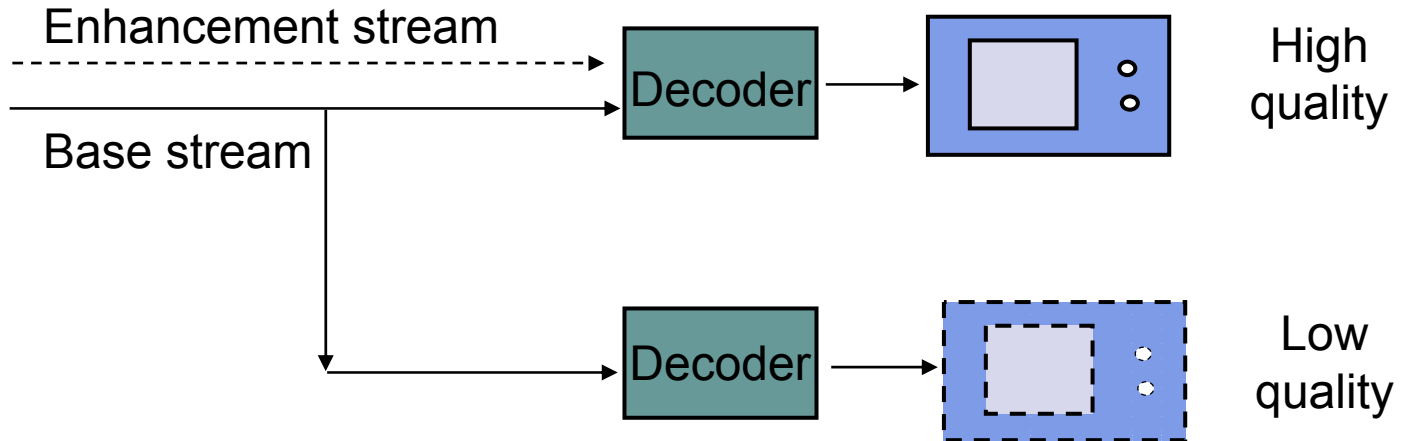
1-D examples:



## Bidirectional Motion Compensation



### SNR scalability



### Spatial Scalability

